

# **IDOS boot documentation**

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# Section 1: IDOS boot protocols

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## 1.1 Sector to iniload protocol

The iniload kernel is loaded to an arbitrary segment. The segment must be at least 60h. Common choices are 60h, 70h, and 200h. At least 1536 bytes of the file must be loaded. Current loaders will load at least 8192 bytes if the file is as large or larger than that. The entrypoint is found by applying no segment adjustment (0) and choosing the offset 400h (1024).

### 1.1.1 File properties

The file must be at least 4096 bytes long. This is now required, beyond the former lower bound of 1536 bytes, to support an optimisation of the FAT12 and FAT16 boot sector loaders. The lDebug loader and the FAT32+FSIBOOT loader currently retain the 1536 bytes limit.

The file may allow multi-use as a flat .COM format executable, flat .SYS format device driver, or MZ .EXE format executable and/or device driver. It is also valid to append arbitrary sized data such as a .ZIP archive.

The file needs to be placed in the root directory for the boot sector loaders. The lDebug loader allows to load a file from any subdirectory and this is also allowed. The file may be fragmented in any part. The file data may be located anywhere in the file system. The supported cluster sizes should be between 32 Bytes and 2 MiB, inclusive. The sector size should be between 32 Bytes and 8 KiB, inclusive.

### 1.1.2 Signatures

At offset 1020 (3FCh) there is the signature '1D'. Behind that there are two bytes with printable non-blank ASCII codepoints. Currently the following signatures are defined:

'1DOS'

IDOS kernel (not yet in use)

'1DRx'

RxDOS kernel

'1DFD'

FreeDOS kernel wrapped in iniload (fdkernpl.asm)

'1Deb'

lDebug

'1DDb'

IDDebug (debuggable IDDebug)

‘lDbC’

ICDebug (conditionally debuggable IDDebug)

‘lDTP’

IDOS test payload kernel (testpl.asm)

‘lDTW’

IDOS test result writer kernel (testwrit.asm)

### 1.1.3 Load Stack Variables (LSV)

Under this protocol, the pointer ‘ss:bp’ is passed. It points to a boot sector with (E)BPB. ‘bp’ must be even for compatibility with older iniload (before 2023-March). The stack pointer must be at most ‘bp - 10h’. Below the pointed to location there live the Load Stack Variables. These follow this structure:

```
        struc LOADSTACKVARS, -10h
lsvFirstCluster:    resd 1
lsvFATSector:       resd 1
lsvFATSeg:          resw 1
lsvLoadSeg:         resw 1
lsvDataStart:       resd 1
        endstruc
```

lsvFirstCluster

(FAT12, FAT16) Low word gives starting cluster of file. High word uninitialised.

(FAT32) Dword gives starting cluster of file.

(else) Should be zero.

lsvFATSector

(FAT16) Low word gives loaded sector-in-FAT. -1 if none loaded yet. High word uninitialised.

(FAT32) Dword gives loaded sector-in-FAT. -1 if none loaded yet.

(FAT12, else) Unused.

lsvFATSeg

(FAT16, FAT32) Word gives segment of FAT buffer if word/dword [lsvFATSector] != -1.

(FAT12) Word gives segment of FAT buffer. Zero if none. Otherwise, buffer holds entire FAT data, up to 6 KiB.

lsvLoadSeg

Word points to segment beyond last loaded paragraph. Allows iniload to determine how much of it is already loaded.

lsvDataStart

Dword gives sector-in-partition of first cluster's data.

An LSV extension allows to pass a command line to the kernel. The base pointer must be at least '114h' then. The stack pointer must be at most 'bp - 114h' then. This follows the structure like this:

```
lsvclSignature          equ "CL"
lsvclBufferLength       equ 256

        struc LSVCMDLINE, LOADSTACKVARS - lsvclBufferLength - 4
lsvCommandLine:
.start:      resb lsvclBufferLength
.signature:  resw 1
lsvExtra:    resw 1
        endstruc
```

lsvCommandLine.start

Command line buffer. Contains zero-terminated command line string.

lsvCommandLine.signature

Contains the signature value 'CL' if command line is given.

lsvExtra

Used internally by iniload. Space for this must be reserved when passing a command line.

If no command line is passed then either the stack pointer must be 'bp - 10h', or 'bp - 12h', or the word in the lsvCommandLine.signature variable (word [ss:bp - 14h]) must not equal the string 'CL'.

- dosemu2's RxDOS.3 support sets 'sp = bp - 10h'
- ldosboot boot.asm (FAT12/FAT16) loader makes sure not to pass the variable with the content "CL". Refer to placeholder and DIRSEARCHSTACK\_CL\_FIRST uses in the source.
- ldosboot boot32.asm (FAT32) loader uses the variable for an 'entries per sector' value which is always a power of two and always below-or-equal 100h.
- lDebug with protocol options cmdline=0 push\_dpt=0 sets 'sp = bp - 10h'

#### 1.1.4 Memory map

The initial loader part that is loaded must be loaded at above or equal to linear 00600h. The FAT buffer segment (if used) must also be stored at above or equal to linear 00600h. The stack (which should extend at least 512 bytes below 'ss:bp') and boot sector (pointed to by 'ss:bp', at least 512 bytes length) should also be stored at above or equal to linear 00600h.

There is an additional memory area, the Low Memory Area top reservation, which should be unused by the load protocol at handoff time but be at least 20 KiB in size. It is located below the usable Low Memory Area top. That is, directly below the EBDA, RPL-reserved memory, video

memory, or otherwise UMA. This area is reserved in order to facilitate initial loader operation.

None of the memory areas may overlap. This does not include the FAT buffer in case it is uninitialised.

### 1.1.5 Load filename in the boot sector

The boot sector may be expected to contain a valid 8.3 format (blank-padded FCB) filename in the area of the boot sector starting behind the (E)BPB, extending up to below the boot sector signature word with value AA55h (at offset 510 in the boot sector). This name should not contain blanks other than trailing in the file name portion or trailing in the file extension portion. It should consist of printable ASCII codepoints. That is, byte values between 20h and 7Eh inclusive. It should not consist of eleven times the same byte value. Additional FAT Short File Name restrictions may be assumed.

Although a loader should not depend on this for crucial operation, it may want to detect the kernel name it was presumably loaded from for informational or optional purposes. The canonical implementation of this is currently the function 'findname' in the `testpl.asm` test payload kernel. It is found within the `ldosboot` repo. This handling is based on the function of the same name in the `instsect` application.

### 1.1.6 Query patch support

The `ldosboot` repo includes a patch Script for `lDebug` (.sld) file which allows to patch the initial loader stage. The patches concern handling of the CHS geometry detection, and whether LBA or CHS access is used. There are several legacy patch sites in which `patch.sld` can directly patch the initial loader's code.

However, the preferred way is to find the query patch sequence. It should appear within the first 1536 bytes, that is within the part of the initial loader that must be loaded. This is the sequence:

```
8A5640  mov dl, byte [bp + 40h]
B8xxxy  mov ax, yyxxh
84D2    test dl, dl
7902    jns @F
86C4    xchg al, ah
@@:
```

The immediate word of the `mov ax` instruction is to be patched. The sequence should be scanned for without regard as to what the current contents of this word are.

The following flag values are used:

- 01h Force CHS access, do not detect LBA support with 13.41
- 02h Force LBA access, do not detect LBA support with 13.41
- 04h Force use of BPB's CHS geometry, do not detect with 13.08
- 80h Used by `lDebug`. If this value is set for the load unit, then `lDebug` will make use of the other flags set up for that unit. The corresponding flags will be saved in `lDebug`'s `load_unit_flags`. This affects only the load unit (LD in `lDebug` terminology), which suffices to pass commands in the startup Script for `lDebug`.

The flag 01h takes precedence over 02h if both are set.

The low byte (xxh) is used in case the loader loads from a diskette unit, that is a unit number below 80h. The high byte (yyh) is used otherwise, in case the loader loads from a hard disk unit, that is a unit number above-or-equal 80h.

## 1.2 Iniload to payload protocol

The payload is loaded to an arbitrary segment. The segment must be at least 60h. The entire payload must be loaded. The size of the payload is determined at iniload build time. The entrypoint is found by applying a segment adjustment and choosing an offset. The segment adjustment is specified at iniload build time by the numeric define `_EXEC_SEGMENT` (default 0), and the offset by the define `_EXEC_OFFSET` (default 0).

### 1.2.1 Stack

The stack must allocate at least 256 bytes of space below `ss : sp`. This space must not underflow, ie `sp` is at least 256 (100h). `sp` should not be higher than 8192 (2000h). `bp` is typically equal to `sp + 120h`. `bp` must be below or equal 7000h. The stack starting at `ss : 0` up to below `ss : bp + 21Ch` should be in the area between the paragraphs indicated by the load top lower limit and memory top upper limit. (Refer to `LOADDATA` fields `ldLoadTop` and `ldMemoryTop`.) Subsequent stages like `inicompl`, `fdkernpl`, `nullpl`, and kernels like `lDebug` and `testpl` may depend on these characteristics.

### 1.2.2 Extended BIO Parameter Block (EBPB)

Above the LSV, `ss : bp` points to an EBPB and surrounding boot sector. Note that this is always a FAT32-style EBPB. If the filesystem that is loaded from is not FAT32, and is therefore FAT16 or FAT12, then the FAT16/FAT12 BPBN structure is moved up. It is placed where the FAT32 BPBN is usually expected. In this case, the entire boot sector contents behind the BPBN are also moved up by the size of the FAT32-specific fields. The FAT32-specific fields are filled with zeros, except for the FAT32 'sectors per FAT' field. It is filled with the contents of the FAT16/FAT12 'sectors per FAT' field.

### 1.2.3 Load Stack Variables (LSV)

Refer to section 1.1.3.

### 1.2.4 Load Data 1 (LD)

Below the LSV, iniload passes the `LOADDATA` (1) structure.

```
        struc LOADDATA, LOADSTACKVARS - 10h
ldMemoryTop:    resw 1
ldLoadTop:      resw 1
ldSectorSeg:    resw 1
ldFATType:      resb 1
ldHasLBA:       resb 1
ldClusterSize: resw 1
ldParaPerSector:resw 1
ldLoadingSeg:
ldQueryPatchValue:
                resw 1
ldLoadUntilSeg: resw 1
        endstruc
```

#### ldMemoryTop

Word. Segment pointer to behind usable memory. Points at the first of the EBDA, RPL-reserved memory, or video memory or otherwise UMA. Indicates how much memory may be used by a typical kernel. (LDebug detects the EBDA to move that below where it installs.)

#### ldLoadTop

Word. Segment pointer to lowest IDOS boot memory in use. All memory between linear 600h and the segment indicated here is usable by the payload. Only the payload itself is stored in this area. The other buffers, stack, and structures passed by iniload must live above this segment.

#### ldSectorSeg

Word. Segment pointer to an 8 KiB transfer buffer. It is insured that this buffer does not cross a 64 KiB boundary. This may be needed by some disk units. The buffer is not initialised to anything generally.

#### ldFATType

Byte. Indicates length of FAT entry in bits. 12 indicates FAT12, 16 FAT16, 32 FAT32. It is planned to allow zero for non-FAT filesystems.

#### ldHasLBA

Byte. Only least significant bit used. Bit on indicates LBA extensions available for the load disk unit. Bit off indicates LBA extensions not available.

#### ldClusterSize

Word. Contains amount of sectors per cluster. Unlike the byte field for the same purpose in the BPB, this field can encode 256 (EDR-DOS compatible) without any masking. May be given as zero for non-FAT filesystems.

#### ldParaPerSector

Word. Contains amount of paragraphs per sector. Must be a power of two between 2 (32 B/s) and 200h (8192 B/s). May be given as zero for non-FAT filesystems.

#### ldLoadingSeg

Word. Internally used by iniload. Available for re-use by payload. However, ldQueryPatchValue re-uses the same field.

#### ldQueryPatchValue

Word. Passes the query patch value from the initial loader. This provides an opportunity to patch a well-known site in the initial loader to change its behaviour in some ways. Near the end of its operation, the initial loader passes along this value in this variable for the next stage to use.

#### ldLoadUntilSeg

Word. Internally used by iniload. Available for re-use by payload.



### 1.2.5 Load Command Line (LCL)

Below the LOADDATA structure, iniload passes the LOADCMDLINE structure.

```
lsvclBufferLength      equ 256

        struc LOADCMDLINE, LOADDATA - lsvclBufferLength
ldCommandLine:
.start:      resb lsvclBufferLength
        endstruc
```

This buffer is always initialised to an ASCIZ string. At most 255 bytes may be initialised to string data. At most the 256th byte is a zero.

If the first word of the buffer is equal to 0FF00h, that is there is an empty command line the terminator of which is followed by a byte with the value 0FFh, then no command line was passed to iniload. Currently IDebug can pass a command line to iniload when loading with its IDOS, RxDOS.2, RxDOS.3, or FreeDOS protocols. When iniload is loaded as a Multiboot1 or Multiboot2 specification kernel, it is also assumed that a command line can be passed.

## 1.3 Application mode protocol

An iniload payload can be loadable as a DOS application. In this case the iniload MZ executable header will point to the payload. The `_IMAGE_EXE` define to iniload will indicate support for this if it is enabled.

The entire payload will be loaded as a program image just behind the PSP. The entrypoint is determined by the `_IMAGE_EXE_CS` and `_IMAGE_EXE_IP` defines. These default to a `-16:256 + 64` entrypoint. Minus 16 means the code segment will be equal to the PSP, just like for flat .COM style executables. The 256 in the instruction pointer calculation skips past the PSP. Thus, the default `_IMAGE_EXE_IP` will enter the payload at offset 64 counting within the payload.

Likewise `_IMAGE_EXE_SS` and `_IMAGE_EXE_SP` specify the initial stack. They default to `-16:0FFFEh`, mimicking flat .COM style load somewhat. (However, iniload doesn't get a chance to push a zero word onto this stack before the first payload is run. Adding a nullpl, checkpl, or inicompl stage can be used to provide this.)

`_IMAGE_EXE_AUTO_STACK` defaults to zero. If it is nonzero it indicates to use the automatic stack placement. Its value, which defaults to 2048 if defined as empty, determines how large the stack should be (at least). The explicit `_IMAGE_EXE_SS` and `_IMAGE_EXE_SP` are ignored if the automatic stack placement is enabled. The automatic stack is placed into its own segment, stretching from `ss:0` to `ss:sp`. The stack pointer will match the value of `_IMAGE_EXE_AUTO_STACK`.

`_IMAGE_EXE_MIN` indicates how much memory to allocate to the process at least. It defaults to 65536 (64 KiB). It is actually used in the `_IMAGE_EXE_MIN_CALC` define to calculate the minimum allocation value. `_IMAGE_EXE_MIN_CALC` is not a numdef, but can be overridden from the assembler command line. This is the default content of this define:

```
%define _IMAGE_EXE_MIN_CALC \
((( _IMAGE_EXE_MIN \
- (payload.actual_end - payload) \
```

```
- 256 \  
+ _IMAGE_EXE_AUTO_STACK) + 15) & ~15)
```

`_IMAGE_EXE_MAX` is a define that defaults to `0FFFFh`. If it is nonzero it is used directly as the maximum allocation value. `0FFFFh` indicates to allocate the largest memory block to the process, just like for flat .COM style load. If it is zero, then the maximum allocation field is set to equal the minimum allocation field.

Most of the stages of iniload payloads can pass through the execution control flow from being loaded as a payload to hand off to their payload. This includes the `nullpl`, `checkpl`, and `inicompl` stages. Each stage must be configured appropriately, enabling the `_IMAGE_EXE` define and setting up the other defines appropriately. These three stages expect `-16:256 + 64` or `-16:256` entrypoints as input for application mode, but their output entrypoints can be configured using the defines listed above.

The registers `ds`, `es`, and `ax` are passed through from `iniload`, `nullpl`, `checkpl`, or `inicompl` to each one's payload. The two segment registers will point to the PSP. The `ax` register contains information on the validity of the drives of the two default FCBs.

The iniload MZ executable header's relocation table is currently always empty. The application has to implement relocations by itself.

`_SECOND_PAYLOAD_EXE` allows to include a different payload that is to be used as the application mode program image. In this case the main iniload payload (`_PAYLOAD_FILE`) is used only for boot loaded mode. A different file can be specified with `_SECOND_PAYLOAD_FILE`. Each `_IMAGE_EXE_*` define has an equivalent define named according to the `_SECOND_PAYLOAD_EXE_*` scheme.

## 1.4 Device mode protocol

The device mode load of the payload shares its program image with the application mode. The very beginning of the image must contain a device header. (The first device header, if more than one exist.)

The `inicompl`, `nullpl`, and `checkpl` stages expect certain offsets for the entrypoints: 18 for the strategy, and 22 for the interrupt. Each passthrough stage must be configured with the appropriate device name and device attributes contents, and with the `_DEVICE` define enabled. (These defines are not needed for `iniload`. However, `_IMAGE_EXE` is required.)

The `inicompl` stage uses the request header's break address as passed by DOS to determine how much memory is available to the depacker. This requires MS-DOS version 5 compatibility.

## Source Control Revision ID

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